🌰 Instance-Optimized String Fingerprints

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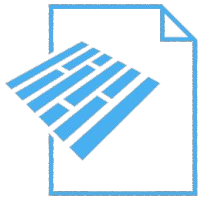
Data Systems Lab x Discrete Optimization Lab @UTN

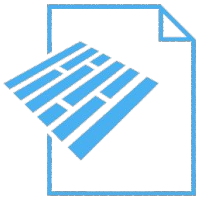
@AIDB'25, September 1st, 2025

UTN Technische Universität Nürnberg

Data Pruning



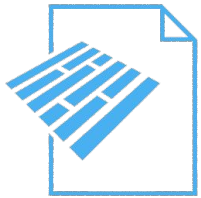
	salary	location
min:	100K	Barcelona
max:	200K	Paris



min:	200K	Madrid
max:	300K	Redmond

Data Pruning

WHERE salary = 250K



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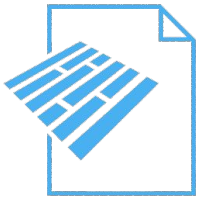
❌ (skip)



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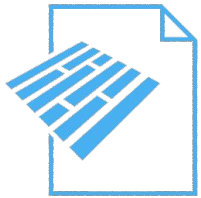
🔍 (maybe)

Data Pruning

❄️ WHERE location LIKE 'B%'



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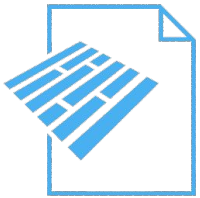
max:	200K	Paris	



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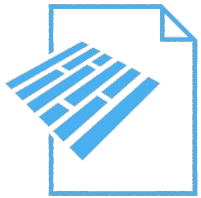
Data Pruning

WHERE location LIKE '%ch%'



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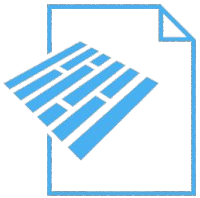

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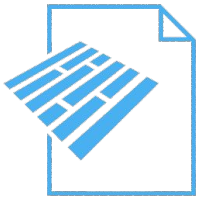

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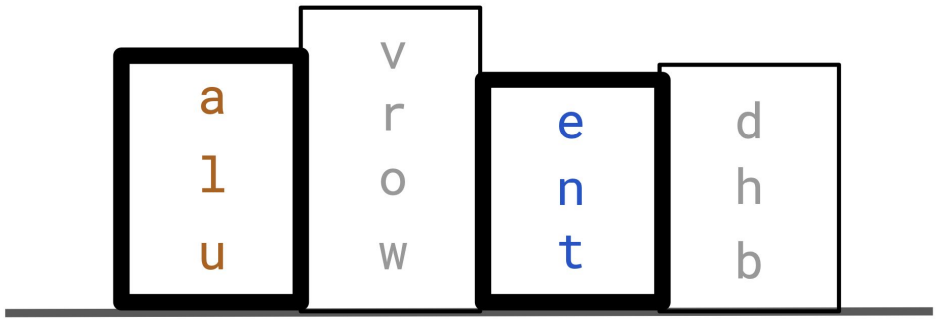

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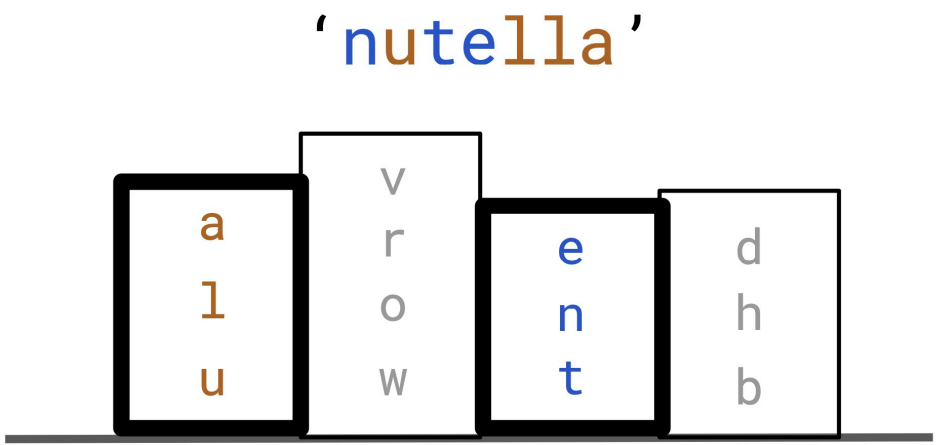

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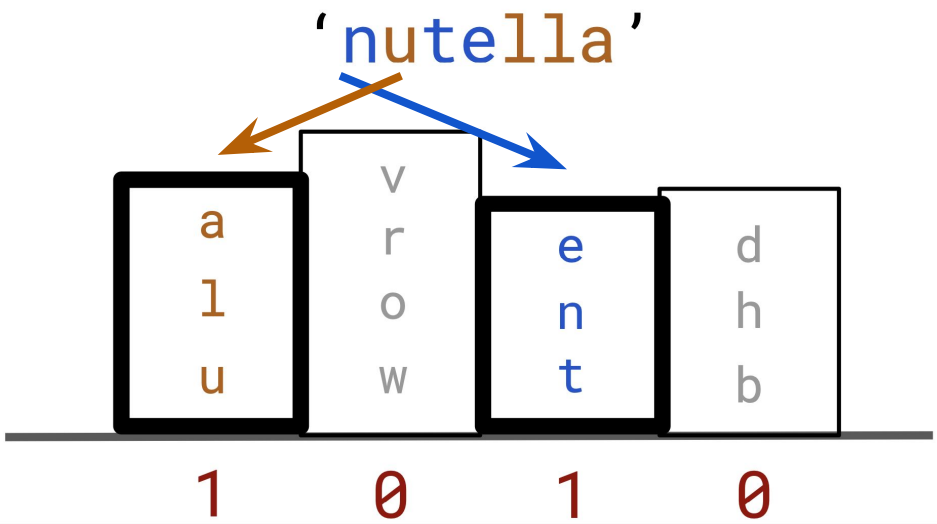

Application

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⇒ Objective: *Minimize the number of false positives.*

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- Words **W**: The string column.
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- Objective: $\max \sum_{q \in Q} \sum_{w \in W \setminus f(q)} \eta_{w,q}$.
- Constraints: Details in the paper.

Evaluation

- Setup: Column `title.title` in IMDb dataset (2.37M tuples; no UTF-8).
- *Workload*:
  - 300 queries ⇒ 10 high-, mid-, low-frequency {1, …, 10}-grams from the column.
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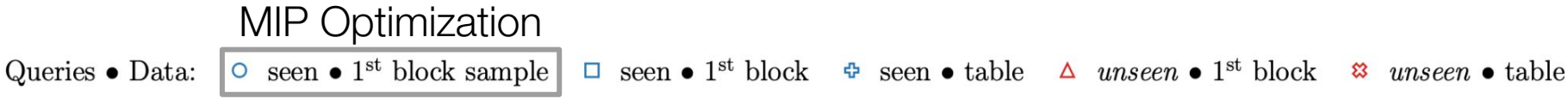
- Baseline: Round-robin placement of letters into bins.

Evaluation: False Positive Rate



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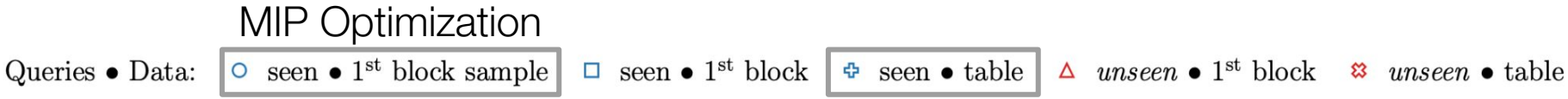

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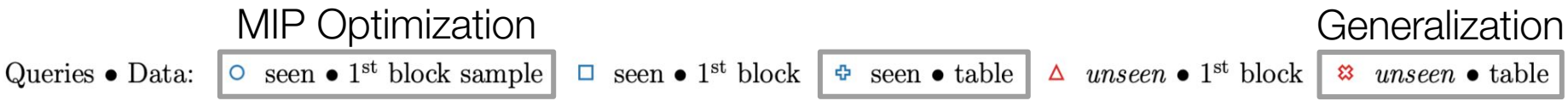

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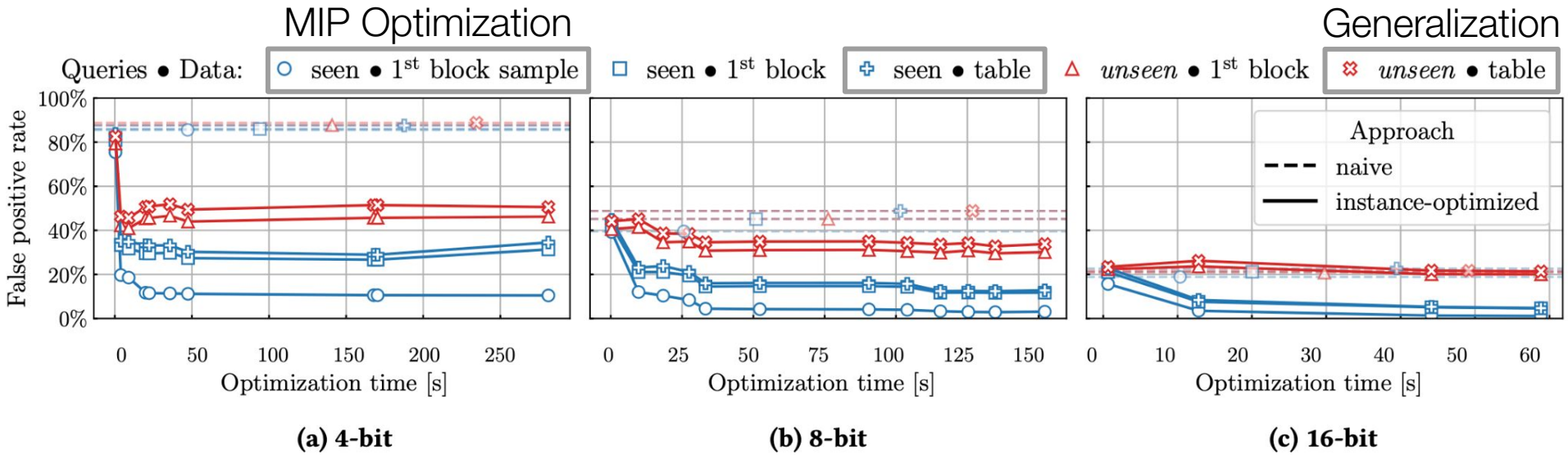
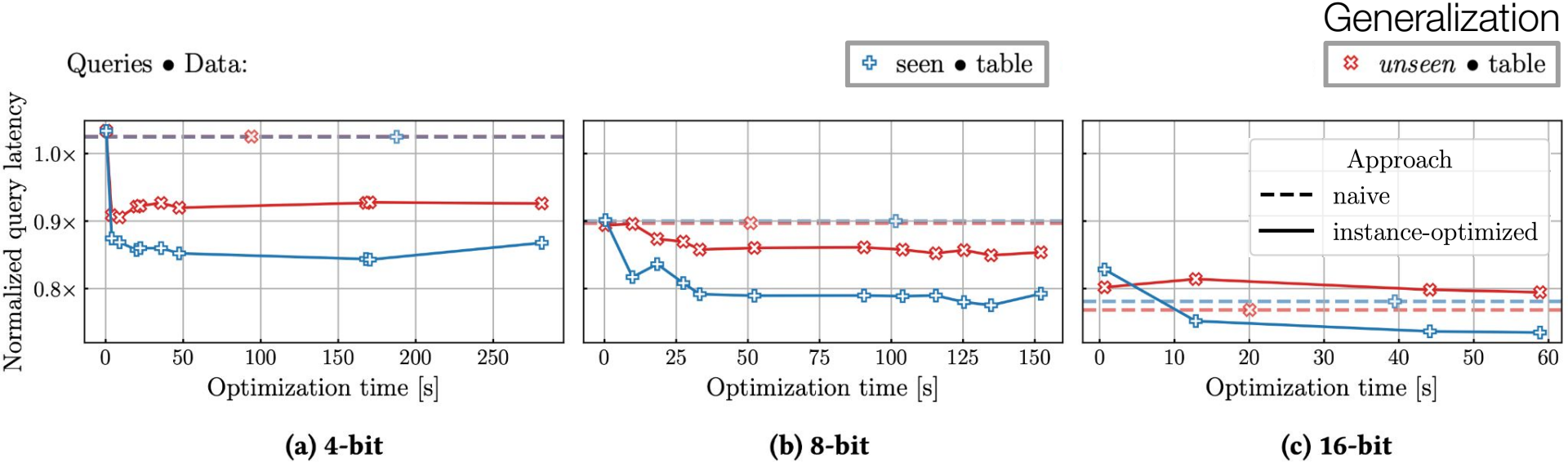


(a) 4-bit **(b) 8-bit** **(c) 16-bit**

Evaluation: Normalized Query Latency

- Note: Run on the *full* table.



Evaluation: Takeaways

- <20% false positive rate on the full table.
- 🪄 Generalization to *unseen* queries (unlike predicate caching).
- 🚀 Up to 1.36x speedup for seen queries & 1.26x speedup for unseen queries.

(Many) Future Work Directions

- Instead of 1-grams, i.e., letters ⇒ Why not 2-/3-grams?
  - Intuition: We can capture the *order* of the letters.
- String zonemaps:
  - String fingerprints enable pruning for infix predicates 😲.
- String cardinality estimation.
  - Take the supersets and sum up their corresponding cardinality.
- Table clustering:
  - Sort by the fingerprint.

Wanna More Cool Research?

See you **tmrw** in *Research 8* (**Westminster 4F**), 1.45pm - 3.15pm!

“Parachute: Single-Pass Bi-Directional Information Passing”

🚀 **1.54x** speedup over `duckdb`